# CSE-101: Discrete Mathematics Chapter 11: Trees 

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## Tree

Definition 1. A tree is a connected undirected graph with no simple circuits.

Theorem 1. An undirected graph is a tree if and only if there is a unique simple path between any two of its vertices.

## Which graphs are trees?

a)

b)


## Specify a vertex as root

Then, direct each edge away from the root.


## Specify a root.

Then, direct each edge away from the root.


## Specify a root.

Then, direct each edge away from the root.


## Specify a root.

Then, direct each edge away from the root.

## ROOT

a)


## What if a different root is chosen?

Then, direct each edge away from the root.


## What if a different root is chosen?

Then, direct each edge away from the root.
a)


## What if a different root is chosen?

Then, direct each edge away from the root.


## What if a different root is chosen?

Then, direct each edge away from the root.


A different rooted tree results.

## Jake's Pizza Shop Tree



## A Tree Has a Root



## A Tree Has a Root



## LEAF NODES

## A Tree Has a Root



## A Tree Has a Root



## A Tree Has a Root



## A Tree Has a Root



## A Tree Has a Root



## A Tree Has a Root



## A Tree Has a Root



RIGHT SUBTREE OF ROOT

## Internal Vertex

A vertex that has children is called an internal vertex.

The subtree at vertex $v$ is the subgraph of the tree consisting of vertex $v$ and its descendants and all edges incident to those descendants.

## How many internal vertices?



## Binary Tree

Definition $2^{\prime \prime}$. A rooted tree is called a binary tree if every internal vertex has no more than 2 children.

The tree is called a full binary tree if every internal vertex has exactly 2 children.

## Ordered Binary Tree

Definition $2^{\prime \prime}$ '. An ordered rooted tree is a rooted tree where the children of each internal vertex are ordered.

In an ordered binary tree, the two possible children of a vertex are called the left child and the right child, if they exist.

## Tree Properties

Theorem 2. A tree with $\mathbf{N}$ vertices has $\mathbf{N}-1$ edges.

Theorem 5. There are at most $2^{\mathrm{H}}$ leaves in a binary tree of height $H$.

Corallary. If a binary tree with $L$ leaves is full and balanced, then its height is

$$
\mathrm{H}=\left\lceil\log _{2} \mathrm{~L}\right\rceil
$$

## An Ordered Binary Tree



## Parent

$\square$ The parent of a non-root vertex is the unique vertex u with a directed edge from $u$ to $v$.

## What is the parent of Ed?



## Leaf

## - A vertex is called a leaf if it has no

 children.
## How many leaves?



## Ancestors

- The ancestors of a non-root vertex are all the vertices in the path from root to this vertex.


## How many ancestors of Ken?



## Descendants

- The descendants of vertex vare all the vertices that have v as an ancestor.


## How many descendants of Hal?



## Level

- The level of vertex $v$ in a rooted tree is the length of the unique path from the root to v .


## What is the level of Ted?



## Height

$\square$ The height of a rooted tree is the maximum of the levels of its vertices.

## What is the height?



## Balanced

- A rooted binary tree of height H is called balanced if all its leaves are at levels H or H-1.


## Is this binary tree balanced?



## Searching takes time ...

So the goal in computer programs is to find any stored item efficiently when all stored items are ordered.

A Binary Search Tree can be used to store items in its vertices. It enables efficient searches.

## A Binary Search Tree (BST) is . . .

A special kind of binary tree in which:

1. Each vertex contains a distinct key value,
2. The key values in the tree can be compared using "greater than" and "less than", and
3. The key value of each vertex in the tree is less than every key value in its right subtree, and greater than every key value in its left subtree.

## Shape of a binary search tree . . .

Depends on its key values and their order of insertion. Insert the elements ' $J$ ' ' $E$ ' ' $F$ ' ' $T$ ' ' $A$ ' in that order. The first value to be inserted is put into the root.

```
`J'
```


## Inserting ' $E$ ' into the BST

Thereafter, each value to be inserted begins by comparing itself to the value in the root, moving left it is less, or moving right if it is greater. This continues at each level until it can be inserted as a new leaf.


## Inserting ' $F$ ' into the BST

Begin by comparing ' $F$ ' to the value in the root, moving left it is less, or moving right if it is greater. This continues until it can be inserted as a leaf.


## Inserting 'T' into the BST

Begin by comparing ' $T$ ' to the value in the root, moving left it is less, or moving right if it is greater. This continues until it can be inserted as a leaf.


## Inserting ' $A$ ' into the BST

Begin by comparing ' $A$ ' to the value in the root, moving left it is less, or moving right if it is greater. This continues until it can be inserted as a leaf.


## What binary search tree . . .

is obtained by inserting
the elements ' $A$ ' ' $E$ ' ' $F$ ' ' $J$ ' ' $T$ ' in that order?
' $A$ '

## Binary search tree . . .

obtained by inserting
the elements ' $A$ ' ' $E$ ' ' $F$ ' ' $J$ ' ' $T$ ' in that order.


## Another binary search tree



Add nodes containing these values in this order:
'D' ‘B' 'L' 'Q' 'S' 'V' 'Z'

## Is ' $F$ ' in the binary search tree?



## Traversal Algorithms

- A traversal algorithm is a procedure for systematically visiting every vertex of an ordered binary tree.
- Tree traversals are defined recursively.
- Three traversals are named:
preorder,
inorder,
postorder.


## PREORDER Traversal Algorithm

Let T be an ordered binary tree with root r .

If T has only $r$, then $r$ is the preorder traversal.
Otherwise, suppose $T_{1}, T_{2}$ are the left and right subtrees at r . The preorder traversal begins by visiting $r$. Then traverses $T_{1}$ in preorder, then traverses $T_{2}$ in preorder.

## Preorder Traversal: J E A H T M Y



## INORDER Traversal Algorithm

Let $T$ be an ordered binary tree with root r .

If $\mathbf{T}$ has only $r$, then $r$ is the inorder traversal.
Otherwise, suppose $T_{1}, T_{2}$ are the left and right subtrees at r. The inorder traversal begins by traversing $\mathrm{T}_{1}$ in inorder. Then visits $r$, then traverses $T_{2}$ in inorder.

## Inorder Traversal: A E H J M T Y



## POSTORDER Traversal Algorithm

Let T be an ordered binary tree with root r .

If $T$ has only $r$, then $r$ is the postorder traversal.
Otherwise, suppose $T_{1}, T_{2}$ are the left and right subtrees at $r$. The postorder traversal begins by traversing $\mathrm{T}_{1}$ in postorder. Then traverses $\mathrm{T}_{2}$ in postorder, then ends by visiting r .

## Postorder Traversal: A H E M Y T J



## A Binary Expression Tree



INORDER TRAVERSAL: 8 - 5 has value 3 PREORDER TRAVERSAL: - 85

POSTORDER TRAVERSAL: 85 -

## A Binary Expression Tree is . . .

A special kind of binary tree in which:

1. Each leaf node contains a single operand,
2. Each nonleaf node contains a single binary operator, and
3. The left and right subtrees of an operator node represent subexpressions that must be evaluated before applying the operator at the root of the subtree.

## A Binary Expression Tree



What value does it have?
$(4+2) * 3=18$

## A Binary Expression Tree



What infix, prefix, postfix expressions does it represent?

## A Binary Expression Tree



| Infix: | $((4+2) * 3)$ |  |
| :--- | :--- | :--- |
| Prefix: | $*+423$ | evaluate from right |
| Postfix: | $42+3$ * | evaluate from left |

## Levels Indicate Precedence

When a binary expression tree is used to represent an expression, the levels of the nodes in the tree indicate their relative precedence of evaluation.

Operations at higher levels of the tree are evaluated later than those below them. The operation at the root is always the last operation performed.

## Evaluate

## this binary expression tree



What infix, prefix, postfix expressions does it represent?

## A binary expression tree



Infix:

$$
((8-5) *((4+2) / 3))
$$

Prefix: *-85/+423
Postfix:
85-42+3/* has operators in order used

## A binary expression tree



Infix:
$((8-5) *((4+2) / 3))$
Prefix:
*-85/+423
evaluate from right
Postfix:
85-42+3/* evaluate from left

## Inorder Traversal: ( $\mathbf{A}+\mathrm{H}$ ) / ( $\mathbf{M} \mathbf{- Y}$ )



Print left subtree first

## Preorder Traversal: / + A H - M Y



Print left subtree second
Print right subtree last

## Postorder Traversal: A H + M Y - /



Print left subtree first
Print right subtree second

## The End

