



ICS141: Discrete Mathematics for Computer Science I

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Lecture 20

Chapter 4. Induction and Recursion

4.1 Mathematical Induction

4.2 Strong Induction



Mathematical Induction

- A powerful, rigorous technique for proving that a statement $P(n)$ is true for **every** positive integers n , no matter how large.
- Essentially a “domino effect” principle.
- Based on a predicate-logic inference rule:

$$P(1)$$

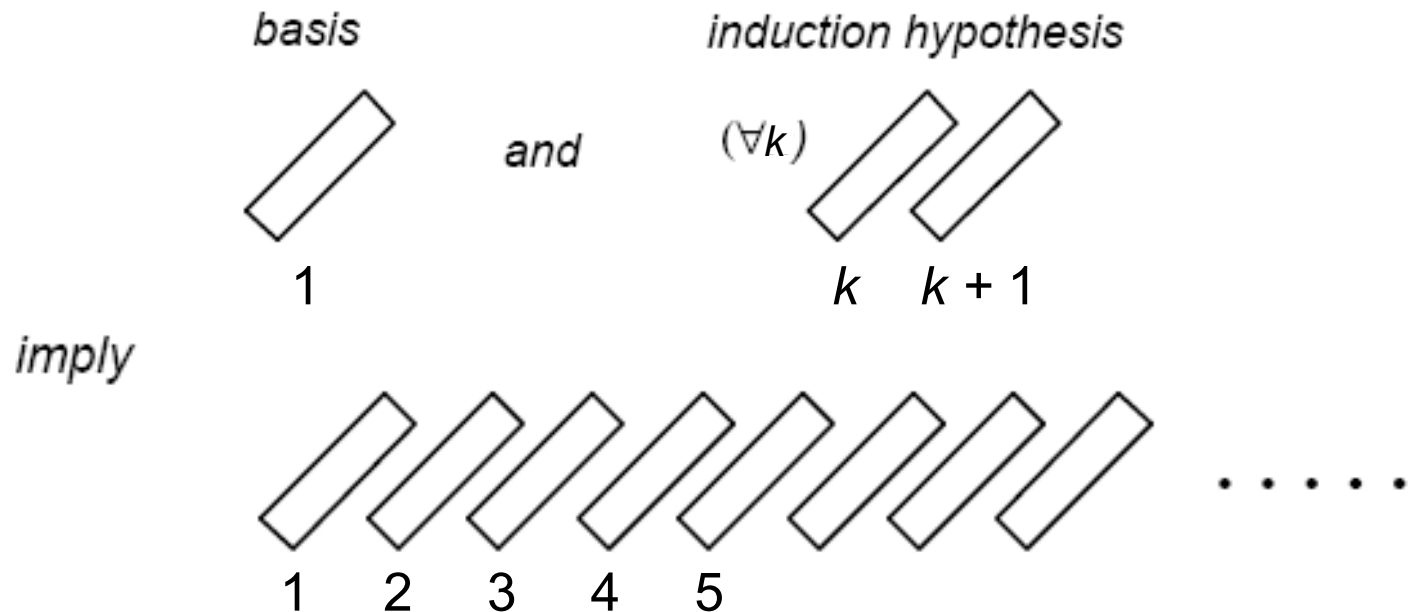
$$\forall k \geq 1 [P(k) \rightarrow P(k+1)]$$

$$\therefore \forall n \geq 1 P(n)$$

*“The First Principle
of Mathematical
Induction”*

The “Domino Effect”

- **Premise #1:** Domino #1 falls.
- **Premise #2:** For every $k \in \mathbb{Z}^+$, if domino # k falls, then so does domino # $k+1$.
- **Conclusion:** All of the dominoes fall down!



Note: this works even if there are infinitely many dominoes!

Mathematical Induction Recap.

- **PRINCIPLE OF MATHEMATICAL INDUCTION:**

To prove that a statement $P(n)$ is true for all positive integers n , we complete two steps:

- **BASIS STEP:** Verify that $P(1)$ is true
- **INDUCTIVE STEP:** Show that the conditional statement $P(k) \rightarrow P(k+1)$ is true for all positive integers k

Inductive Hypothesis

Validity of Induction

Proof: that $\forall n \geq 1 P(n)$ is a valid consequent:

Given any $k \geq 1$, the 2nd premise

$\forall k \geq 1 (P(k) \rightarrow P(k+1))$ trivially implies that

$(P(1) \rightarrow P(2)) \wedge (P(2) \rightarrow P(3)) \wedge \dots \wedge (P(n-1) \rightarrow P(n))$.

Repeatedly applying the hypothetical syllogism rule to adjacent implications in this list $n - 1$ times then gives us $P(1) \rightarrow P(n)$; which together with $P(1)$ (premise #1) and *modus ponens* gives us $P(n)$.

Thus $\forall n \geq 1 P(n)$. ■

Outline of an Inductive Proof

- Let us say we want to prove $\forall n \in \mathbb{Z}^+ P(n)$.
 - Do the **base case** (or **basis step**):
Prove $P(1)$.
 - Do the **inductive step**:
Prove $\forall k \in \mathbb{Z}^+ P(k) \rightarrow P(k+1)$.
 - E.g. you could use a direct proof, as follows:
 - Let $k \in \mathbb{Z}^+$, assume $P(k)$. (*inductive hypothesis*)
 - Now, under this assumption, prove $P(k+1)$.
 - The inductive inference rule then gives us $\forall n \in \mathbb{Z}^+ P(n)$.

Induction Example

- Show that, for $n \geq 1$

$$1 + 2 + \cdots + n = \frac{n(n+1)}{2}$$

- Proof by induction

- $P(n)$: the sum of the first n positive integers is $n(n+1)/2$, i.e. $P(n)$ is
- **Basis step**: Let $n = 1$. The sum of the first positive integer is 1, i.e. $P(1)$ is true.

$$1 = \frac{1(1+1)}{2}$$

Example (cont.)

- **Inductive step:** Prove $\forall k \geq 1: P(k) \rightarrow P(k+1)$.
 - Inductive Hypothesis, $P(k)$:

$$1 + 2 + \cdots + k = \frac{k(k+1)}{2}$$

- Let $k \geq 1$, assume $P(k)$, and prove $P(k+1)$, i.e.

$$\begin{aligned} 1 + 2 + \cdots + k + (k+1) &= \frac{(k+1)[(k+1)+1]}{2} \\ &= \frac{(k+1)(k+2)}{2} \end{aligned}$$

This is what
you have to
prove

$\underbrace{\hspace{1.5cm}}_{P(k+1)}$

Example (cont.)

- **Inductive step** continues...

By inductive hypothesis $P(k)$

$$\begin{aligned} \underbrace{1 + 2 + \cdots + k}_{P(k+1)} + (k+1) &= \frac{k(k+1)}{2} + (k+1) \\ &= \frac{k(k+1)}{2} + \frac{2(k+1)}{2} \\ &= \frac{k^2 + 3k + 2}{2} \\ &= \frac{(k+1)(k+2)}{2} \end{aligned}$$

- Therefore, by the principle of mathematical induction $P(n)$ is true for all integers n with $n \geq 1$



Induction Example 2

- **Example 2:** Conjecture a formula for the sum of the first n positive odd integers. Then prove your conjecture using mathematical induction.

- Practical Method for General Problem Solving.

Special Case: Deriving a Formula

Step 1. Calculate the result for some small cases

Step 2. Guess a formula to match all those cases

Step 3. Verify your guess in the general case

Example 2 (cont.)

- Step 1: Examine small cases

$$1 = 1 = 1^2$$

$$1 + 3 = 4 = 2^2$$

$$1 + 3 + 5 = 9 = 3^2$$

$$1 + 3 + 5 + 7 = 16 = 4^2$$


- Step 2: It sure looks like $1 + 3 + \cdots + (2n-1) = n^2$
- Step 3: Try to prove this assertion by induction

$$\forall n \geq 1, \sum_{i=1}^n (2i-1) = n^2$$

Example *Continues...*

- Prove that the sum of the first n odd positive integers is n^2 . That is, prove:

$$\forall n \geq 1, \sum_{i=1}^n (2i-1) = n^2$$


 $P(n)$

- Proof by induction.
 - **Basis step:** Let $n = 1$. The sum of the first 1 odd positive integer is 1 which equals 1^2 . i.e. $P(1)$ is true.

Example Continues...

- **Inductive step:** Prove $\forall k \geq 1: P(k) \rightarrow P(k+1)$.

- Inductive Hypothesis, $P(k): \left[\sum_{i=1}^k (2i-1) = k^2 \right]$

- Let $k \geq 1$, assume $P(k)$, and prove $P(k+1)$.

$$\begin{aligned} \sum_{i=1}^{k+1} (2i-1) &= \left(\sum_{i=1}^k (2i-1) \right) + (2(k+1)-1) \\ &= k^2 + (2k+1) \quad \text{By inductive hypothesis } P(k) \\ &= (k+1)^2 \end{aligned}$$

Induction Example 3

- Prove that $\forall n \geq 1, n < 2^n$. Let $P(n) = (n < 2^n)$
 - **Basis step:** $P(1): (1 < 2^1) \equiv (1 < 2)$: **True**.
 - **Inductive step:** For $k \geq 1$, prove $P(k) \rightarrow P(k+1)$.
 - Assuming $k < 2^k$, prove $k + 1 < 2^{k+1}$.
 - Note $k + 1 < 2^k + 1$ (by inductive hypothesis)
 $< 2^k + 2^k$ (because $1 < 2^k$ for $k \geq 1$)
 $= 2 \cdot 2^k = 2^{k+1}$
 - So $k + 1 < 2^{k+1}$, i.e. $P(k+1)$ is true
 - Therefore, by the principle of mathematical induction $P(n)$ is true for all integers n with $n \geq 1$.



Generalizing Induction

- Rule can also be used to prove $\forall n \geq c P(n)$ for a given constant $c \in \mathbf{Z}$, where maybe $c \neq 1$.
 - In this circumstance,
the basis step is to prove $P(c)$ rather than $P(1)$,
and the inductive step is to prove $\forall k \geq c (P(k) \rightarrow P(k+1))$.

Induction Example 4

- **Example 6:** Prove that $2^n < n!$ for $n \geq 4$ using mathematical induction.

- $P(n): 2^n < n!$

- **Basis step:** Show that $P(4)$ is true

- Since $2^4 = 16 < 4! = 24$, $P(4)$ is true

- **Inductive step:** Show that $P(k) \rightarrow P(k+1)$ for $k \geq 4$

<div style="border: 1px solid green; padding: 5px; display: inline-block;">$P(k+1)$ is true</div>	{	■ $2^{k+1} = 2 \cdot 2^k$ (by definition of exponent)
		$< 2 \cdot k!$ (by the inductive hypothesis $P(k)$)
		$< (k + 1) \cdot k!$ (because $2 < k+1$ for $k \geq 4$)
		$= (k + 1)!$ (by definition of factorial function)

- Therefore, by the principle of mathematical induction $P(n)$ is true for all integers n with $n \geq 4$.

Second Principle of Induction

a.k.a. “Strong Induction”

- Characterized by another inference rule:

$$\begin{array}{l} P(1) \quad \underbrace{P \text{ is true in } \textit{all} \text{ previous cases}} \\ \forall k \geq 1: (P(1) \wedge P(2) \wedge \cdots \wedge P(k)) \rightarrow P(k+1) \\ \hline \therefore \forall n \geq 1: P(n) \end{array}$$

- The only difference between this and the 1st principle is that:
 - the inductive step here makes use of the stronger hypothesis that all of $P(1), P(2), \dots, P(k)$ are true, not just $P(k)$.

Example of Second Principle

- Show that every integer $n > 1$ can be written as a product $n = p_1 p_2 \dots p_s = \prod p_i$ of some series of s prime numbers.
 - Let $P(n)$ = “ n has that property”
 - **Basis step:** $n = 2$, let $s = 1$, $p_1 = 2$. Then $n = p_1$
 - **Inductive step:** Let $k \geq 2$. Assume $\forall 2 \leq i \leq k: P(i)$.
 - Consider $k + 1$. If it's prime, let $s = 1$, $p_1 = k + 1$.
 - Else $k + 1 = ab$, where $1 < a \leq k$ and $1 < b \leq k$.
Then $a = p_1 p_2 \dots p_t$ and $b = q_1 q_2 \dots q_u$.
(by Inductive Hypothesis)
- Then we have that $k + 1 = p_1 p_2 \dots p_t q_1 q_2 \dots q_u$,
a product of $s = t + u$ primes.

Generalizing Strong Induction

- Handle cases where the inductive step is valid only for integers greater than a particular integer
 - $P(n)$ is true for $\forall n \geq b$ (b : fixed integer)
- **BASIS STEP**: Verify that $P(b), P(b+1), \dots, P(b+j)$ are true (j : a fixed positive integer)
- **INDUCTIVE STEP**: Show that the conditional statement $[P(b) \wedge P(b+1) \wedge \dots \wedge P(k)] \rightarrow P(k+1)$ is true for all positive integers $k \geq b + j$



Another 2nd Principle Example

- Prove that every amount of postage of 12 cents or more can be formed using just 4-cent and 5-cent stamps.
 - $P(n)$ = “postage of n cents can be formed using 4-cent and 5-cent stamps” for $n \geq 12$.
 - ***Basis step:***
 - $12 = 3 \cdot 4$
 - $13 = 2 \cdot 4 + 1 \cdot 5$
 - $14 = 1 \cdot 4 + 2 \cdot 5$
 - $15 = 3 \cdot 5$
 - So $\forall 12 \leq i \leq 15, P(i)$.

Example (cont.)

- **Inductive step:**

- Let $k \geq 15$, assume $\forall 12 \leq i \leq k, P(i)$.

- Note $12 \leq k - 3 \leq k$, so $P(k - 3)$.

(by inductive hypothesis)

This means we can form postage of $k - 3$ cents using just 4-cent and 5-cent stamps.

- Add a 4-cent stamp to get postage for $k + 1$, i.e. $P(k + 1)$ is true (postage of $k + 1$ cents can be formed using 4-cent and 5-cent stamps).

- Therefore, by the 2nd principle of mathematical induction $P(n)$ is true for all integers n with $n \geq 12$.



Another 2nd Principle Example

- Prove by the 1st Principle.
 - $P(n)$ = “postage of n cents can be formed using 4-cent and 5-cent stamps”, $n \geq 12$.
 - **Basis step:** $P(12)$: $12 = 3 \cdot 4$.
 - **Inductive step:** $P(k) \rightarrow P(k+1)$
 - Case 1: At least one 4-cent stamp was used for $P(k)$
 - $k + 1 = k - 4 + 5$ (i.e. replace the 4-cent stamp with a 5-cent stamp to form a postage of $k + 1$ cents)

Example *Continues...*

- **Inductive step:** $P(k) \rightarrow P(k+1)$
 - Case 2: No 4-cent stamps were used for $P(k)$
 - Since $k \geq 12$, at least three 5-cent stamps are needed to form postage of k cents
 - $k + 1 = k - 3 \cdot 5 + 4 \cdot 4$ (i.e. replace three 5-cent stamps with four 4-cent stamps to form a postage of $k + 1$ cents)
- Therefore, by the principle of mathematical induction $P(n)$ is true for all integers n with $n \geq 12$.

The Well-Ordering Property

- Another way to prove the validity of the inductive inference rule is by using the *well-ordering property*, which says that:
 - Every non-empty set of non-negative integers has a minimum (smallest) element.
 - $\forall \emptyset \subset S \subseteq \mathbb{N}: \exists m \in S$ such that $\forall n \in S, m \leq n$
- This implies that $\{n \mid \neg P(n)\}$ (if non-empty) has a minimum element m , but then the assumption that $P(m-1) \rightarrow P((m-1)+1)$ would be contradicted.